

# Graph Overview (1A)

---

Copyright (c) 2015 - 2018 Young W. Lim.

Permission is granted to copy, distribute and/or modify this document under the terms of the GNU Free Documentation License, Version 1.2 or any later version published by the Free Software Foundation; with no Invariant Sections, no Front-Cover Texts, and no Back-Cover Texts. A copy of the license is included in the section entitled "GNU Free Documentation License".

Please send corrections (or suggestions) to [youngwlim@hotmail.com](mailto:youngwlim@hotmail.com).

This document was produced by using LibreOffice and Octave.

# Some class of graphs (1)

## **Complete graph**

A complete graph is a graph in which each pair of vertices is joined by an edge. A complete graph contains all possible edges.

## **Connected graph**

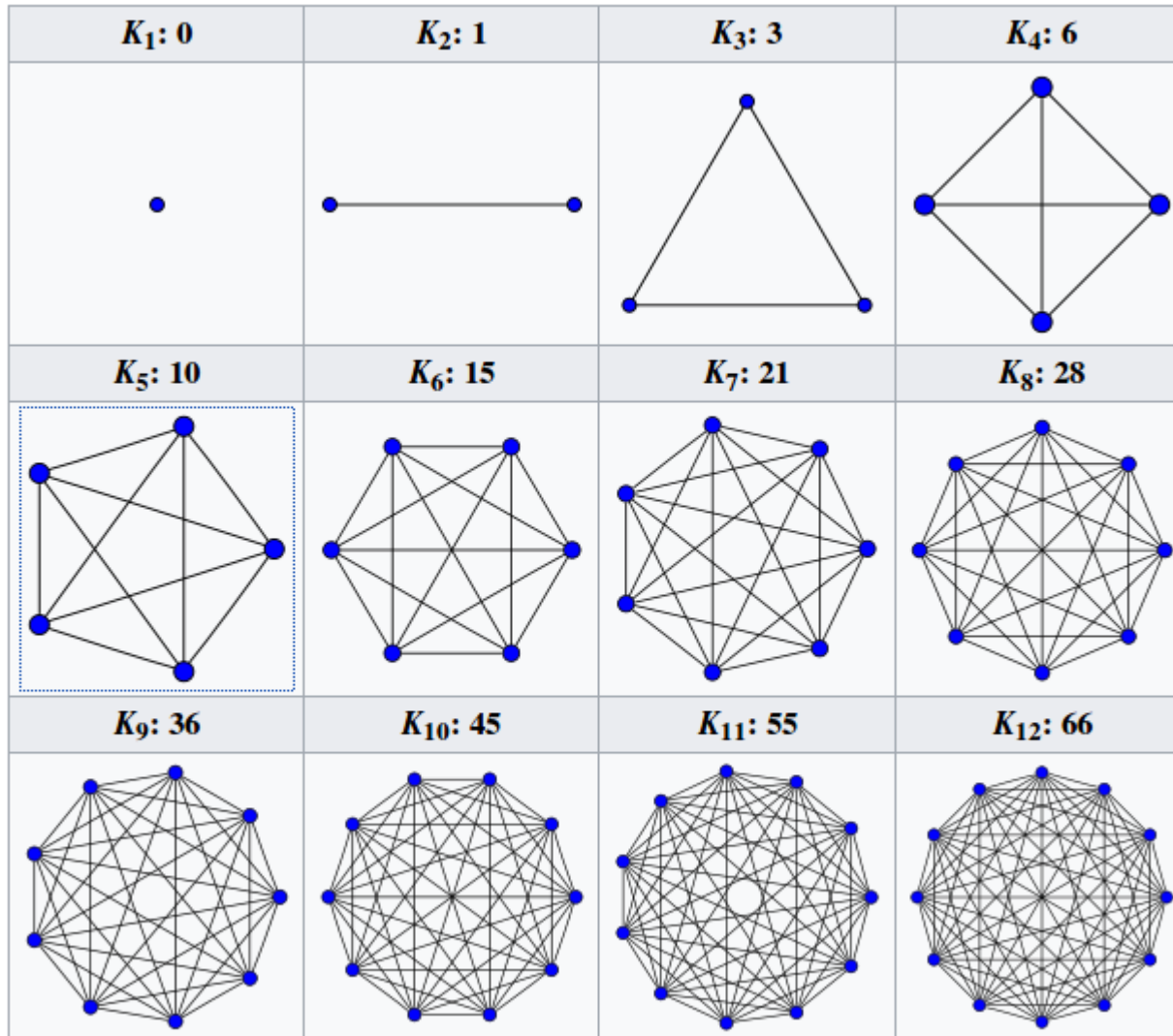
In an undirected graph, an unordered pair of vertices  $\{x, y\}$  is called connected if a path leads from  $x$  to  $y$ . Otherwise, the unordered pair is called disconnected.

## **Bipartite graph**

A bipartite graph is a graph in which the vertex set can be partitioned into two sets,  $W$  and  $X$ , so that no two vertices in  $W$  share a common edge and no two vertices in  $X$  share a common edge. Alternatively, it is a graph with a chromatic number of 2.

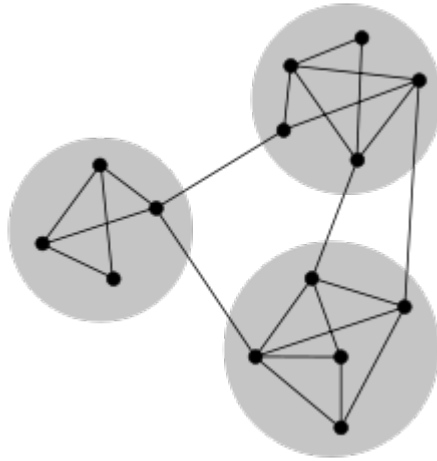
[https://en.wikipedia.org/wiki/Graph\\_\(discrete\\_mathematics\)](https://en.wikipedia.org/wiki/Graph_(discrete_mathematics))

# Complete Graphs

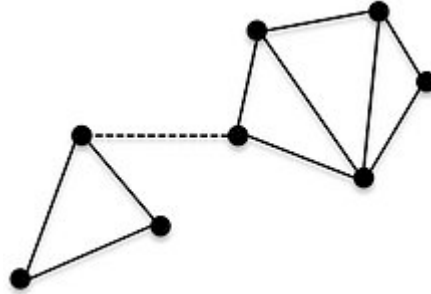


[https://en.wikipedia.org/wiki/Complete\\_graph](https://en.wikipedia.org/wiki/Complete_graph)

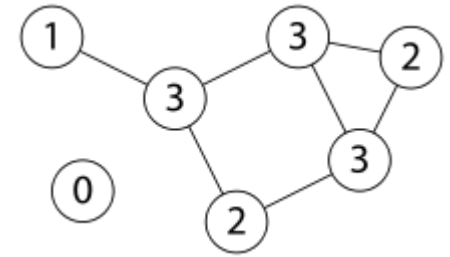
# Connected Graphs



This graph becomes disconnected when the right-most node in the gray area on the left is removed



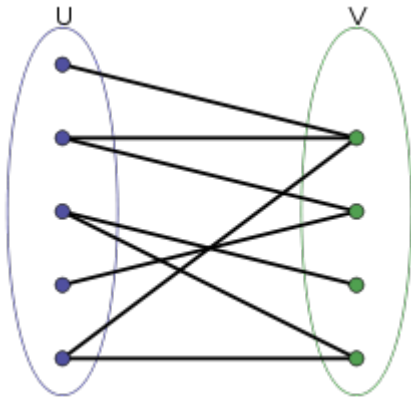
This graph becomes disconnected when the dashed edge is removed.



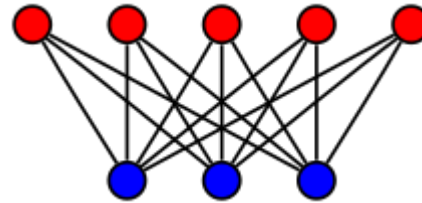
With vertex 0 this graph is disconnected, the rest of the graph is connected.

[https://en.wikipedia.org/wiki/Connectivity\\_\(graph\\_theory\)](https://en.wikipedia.org/wiki/Connectivity_(graph_theory))

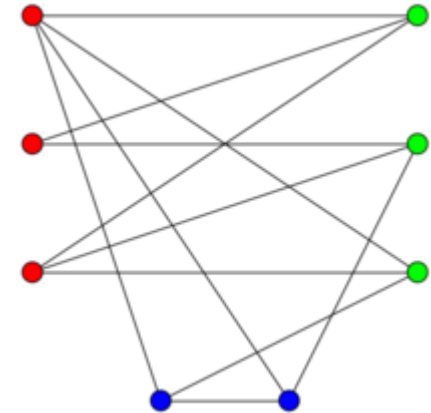
# Bipartite Graphs



Example of a bipartite graph without cycles



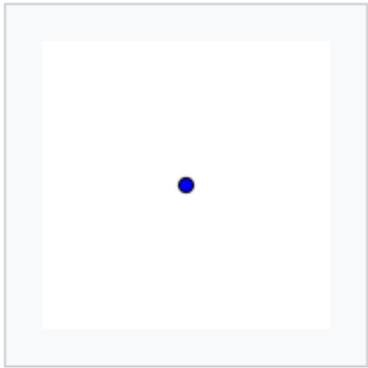
A complete bipartite graph with  $m = 5$  and  $n = 3$



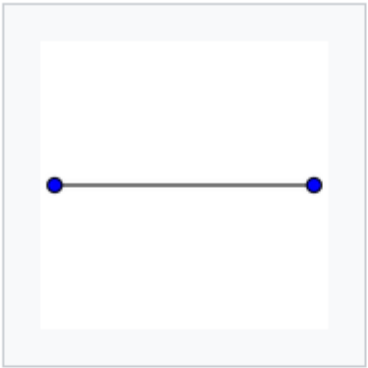
A graph with an odd cycle transversal of size 2: removing the two blue bottom vertices leaves a bipartite graph.

[https://en.wikipedia.org/wiki/Bipartite\\_graph](https://en.wikipedia.org/wiki/Bipartite_graph)

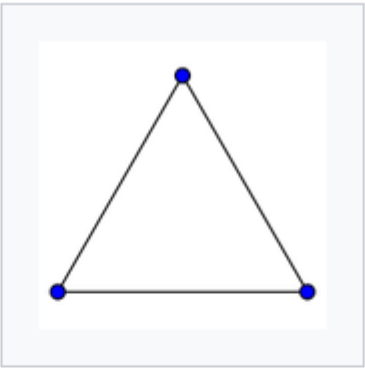
# Complete Graphs



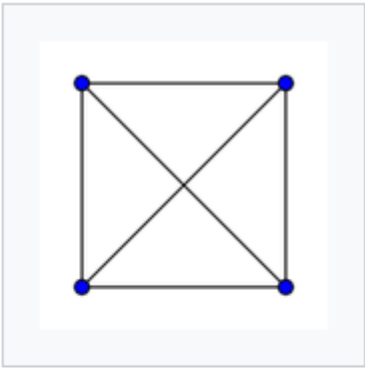
$K_1$



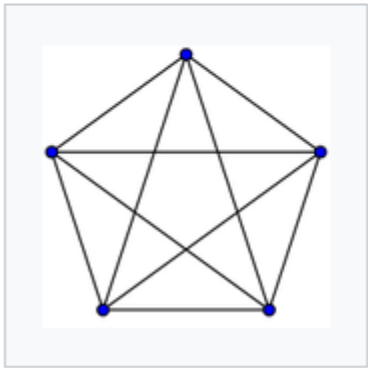
$K_2$



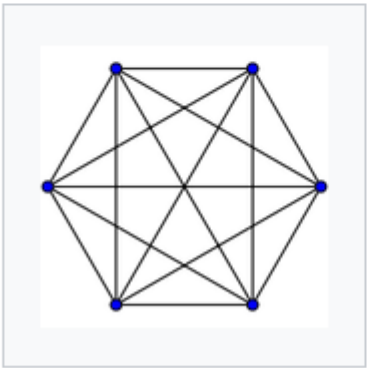
$K_3$



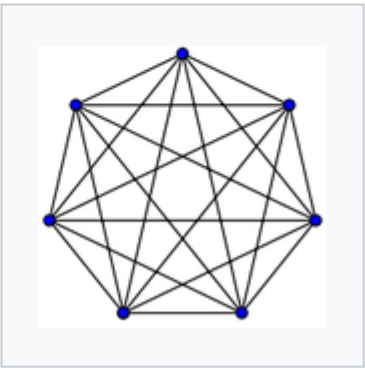
$K_4$



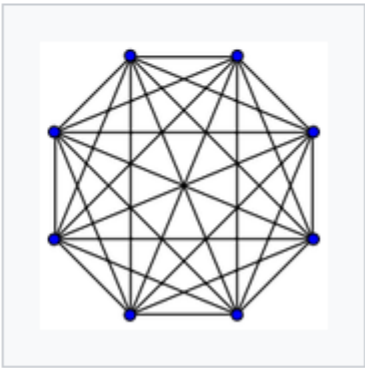
$K_5$



$K_6$



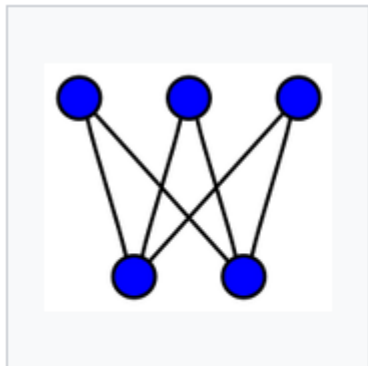
$K_7$



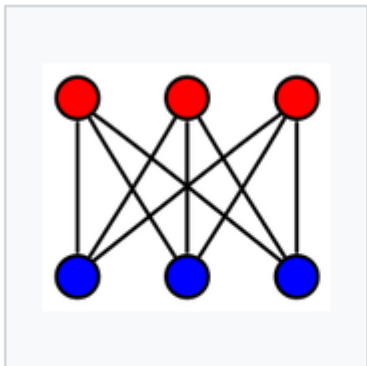
$K_8$

[https://en.wikipedia.org/wiki/Gallery\\_of\\_named\\_graphs](https://en.wikipedia.org/wiki/Gallery_of_named_graphs)

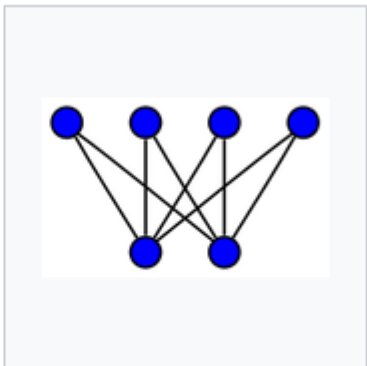
# Complete Bipartite Graphs



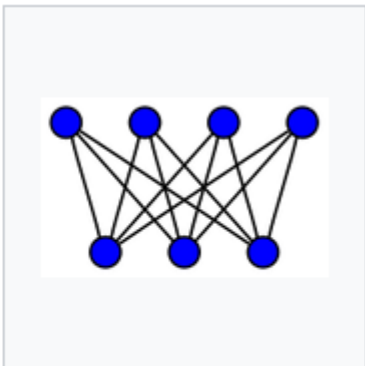
$K_{2,3}$



$K_{3,3}$ , the utility graph



$K_{2,4}$

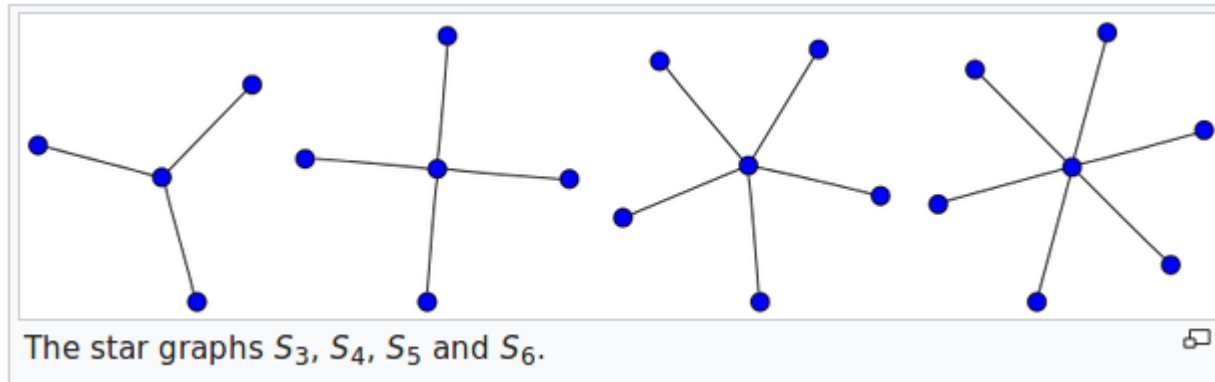


$K_{3,4}$

[https://en.wikipedia.org/wiki/Gallery\\_of\\_named\\_graphs](https://en.wikipedia.org/wiki/Gallery_of_named_graphs)

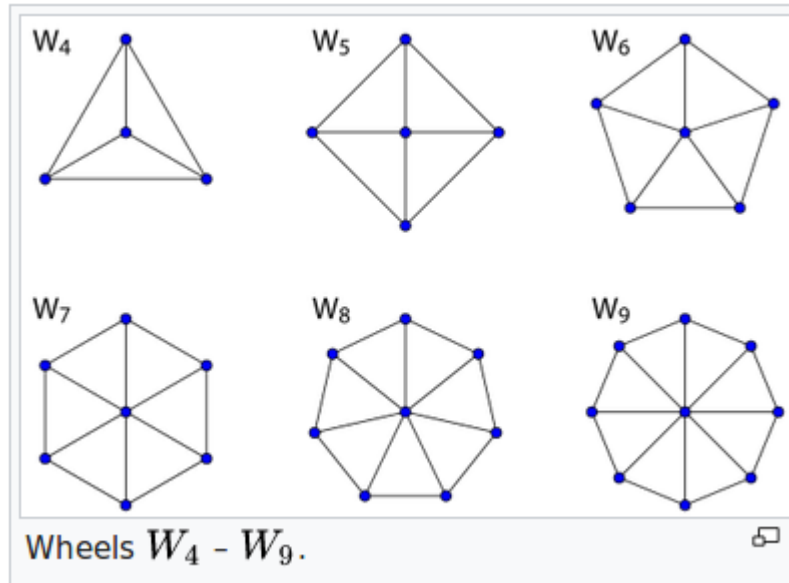


# Star Graphs



[https://en.wikipedia.org/wiki/Gallery\\_of\\_named\\_graphs](https://en.wikipedia.org/wiki/Gallery_of_named_graphs)

# Wheel Graphs



[https://en.wikipedia.org/wiki/Gallery\\_of\\_named\\_graphs](https://en.wikipedia.org/wiki/Gallery_of_named_graphs)

# Some class of graphs (2)

## Planar graph

A planar graph is a graph whose vertices and edges can be drawn in a plane such that no two of the edges intersect.

## Cycle graph

A cycle graph or circular graph of order  $n \geq 3$  is a graph in which the vertices can be listed in an order  $v_1, v_2, \dots, v_n$  such that the edges are the  $\{v_i, v_{i+1}\}$  where  $i = 1, 2, \dots, n - 1$ , plus the edge  $\{v_n, v_1\}$ .

Cycle graphs can be characterized as connected graphs in which the degree of all vertices is 2.


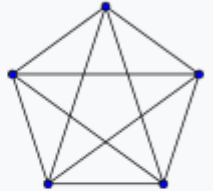

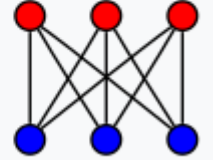
If a cycle graph occurs as a subgraph of another graph, it is a cycle or circuit in that graph.

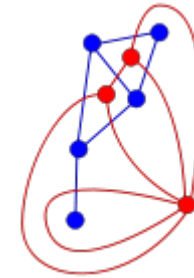
## Tree

A tree is a connected graph with no cycles.

[https://en.wikipedia.org/wiki/Graph\\_\(discrete\\_mathematics\)](https://en.wikipedia.org/wiki/Graph_(discrete_mathematics))

# Planar Graphs

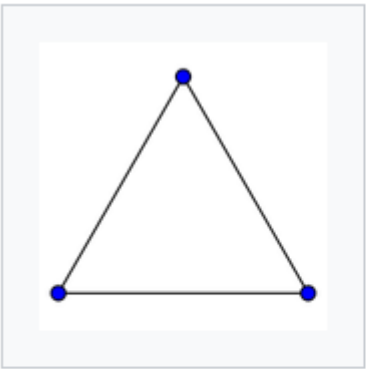
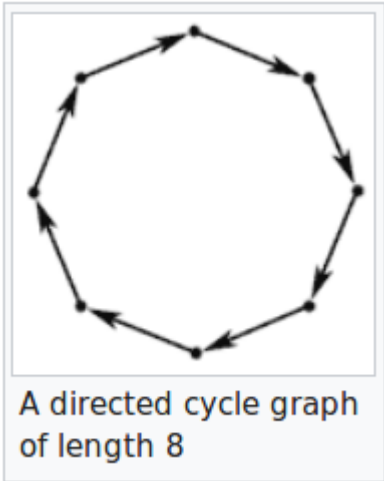
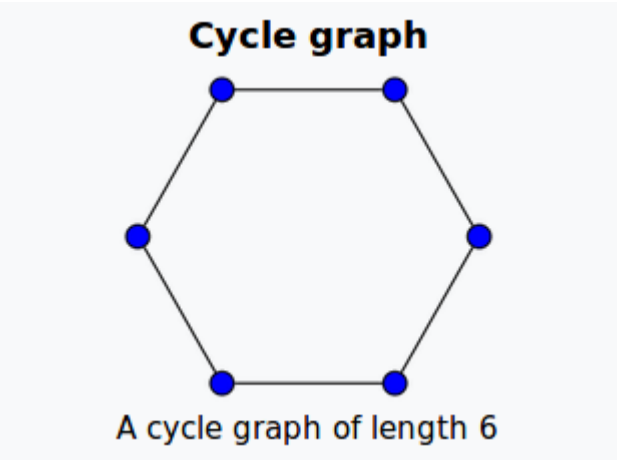
Example graphs	
Planar	Nonplanar
 <p>Butterfly graph</p>	 <p>Complete graph <math>K_5</math></p>
 <p>Complete graph <math>K_4</math></p>	 <p>Utility graph <math>K_{3,3}</math></p>



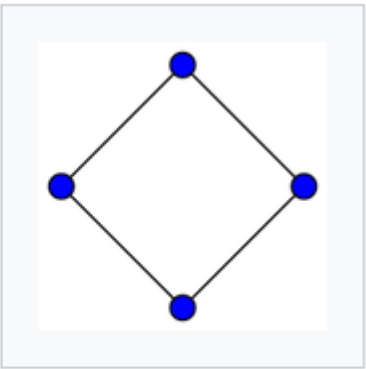
A planar graph and its dual

[https://en.wikipedia.org/wiki/Planar\\_graph](https://en.wikipedia.org/wiki/Planar_graph)

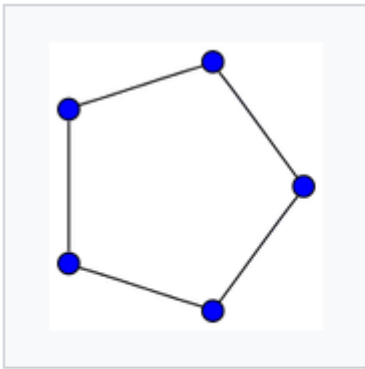
# Cycle Graphs



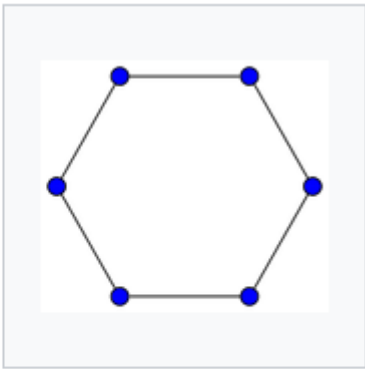
$C_3$



$C_4$



$C_5$

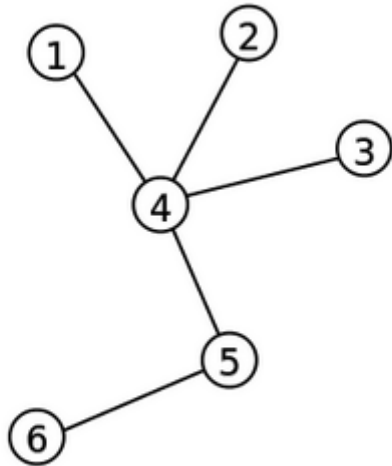


$C_6$

[https://en.wikipedia.org/wiki/Cycle\\_graph](https://en.wikipedia.org/wiki/Cycle_graph)  
[https://en.wikipedia.org/wiki/Gallery\\_of\\_named\\_graphs](https://en.wikipedia.org/wiki/Gallery_of_named_graphs)

# Tree Graphs

**Trees**



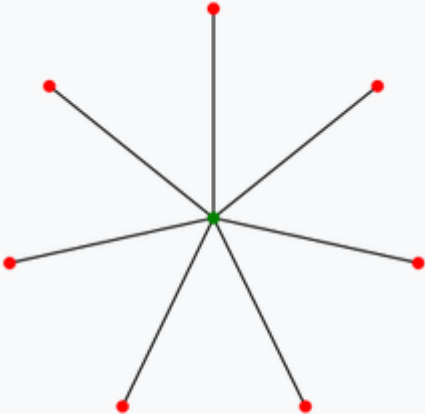
A labeled tree with 6 vertices and 5 edges.

**Path graph**

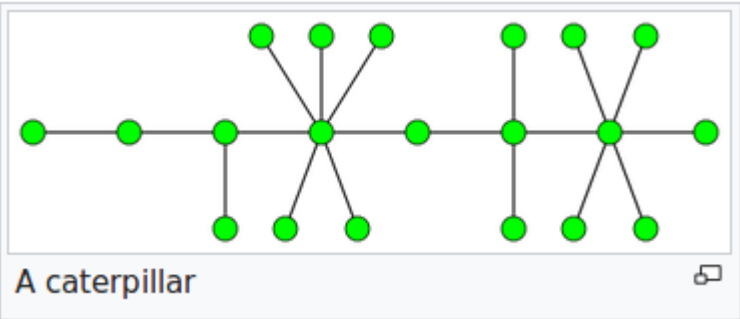


A path graph on 6 vertices

**Star**



The star  $S_7$ . (Some authors index this as  $S_8$ .)



A caterpillar

[https://en.wikipedia.org/wiki/Cycle\\_graph](https://en.wikipedia.org/wiki/Cycle_graph)

# Hypercube

A hypercube can be defined by increasing the numbers of dimensions of a shape:

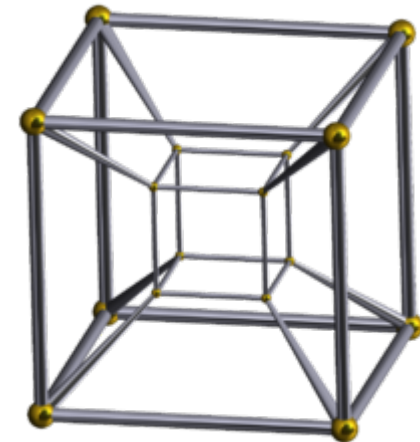
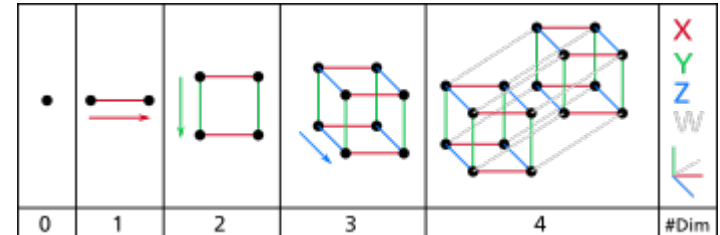
0 – A point is a hypercube of dimension zero.

1 – If one moves this point one unit length, it will sweep out a line segment, which is a unit hypercube of dimension one.

2 – If one moves this line segment its length in a perpendicular direction from itself; it sweeps out a 2-dimensional square.

3 – If one moves the square one unit length in the direction perpendicular to the plane it lies on, it will generate a 3-dimensional cube.

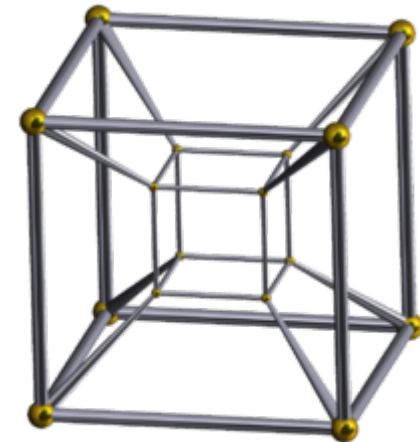
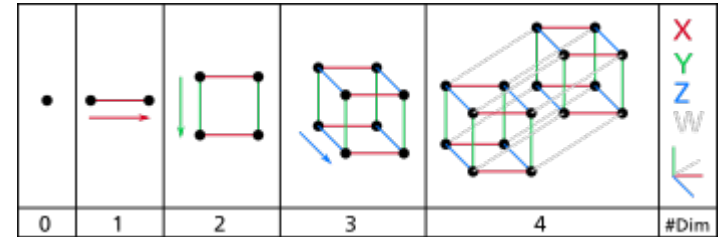
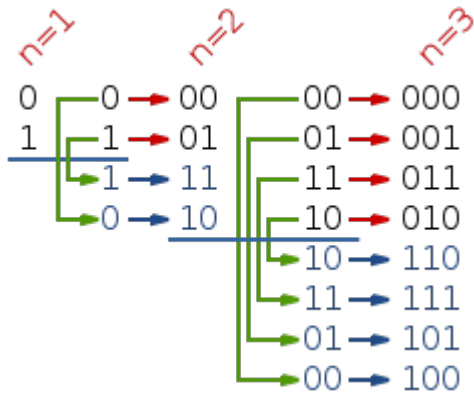
4 – If one moves the cube one unit length into the fourth dimension, it generates a 4-dimensional unit hypercube (a unit tesseract).



**Tesseract**

<https://en.wikipedia.org/wiki/Hypercube>

# Gray Code

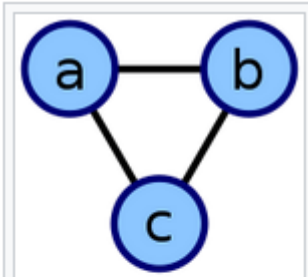


**Tesseract**

[https://en.wikipedia.org/wiki/Gray\\_code](https://en.wikipedia.org/wiki/Gray_code)



# Adjacency Lists



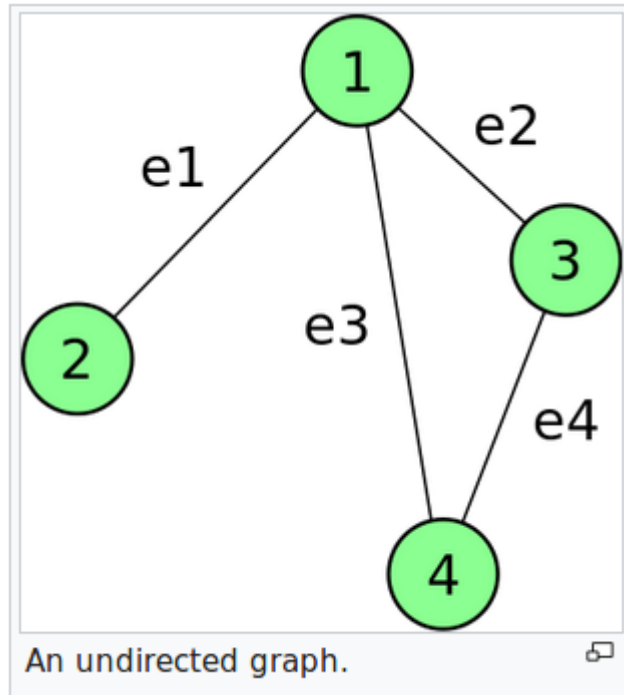
This undirected cyclic graph can be described by the three unordered lists {b, c}, {a, c}, {a, b}.

The graph pictured above has this adjacency list representation:

a	adjacent to	b, c
b	adjacent to	a, c
c	adjacent to	a, b

[https://en.wikipedia.org/wiki/Adjacency\\_list](https://en.wikipedia.org/wiki/Adjacency_list)

# Incidence Matrix

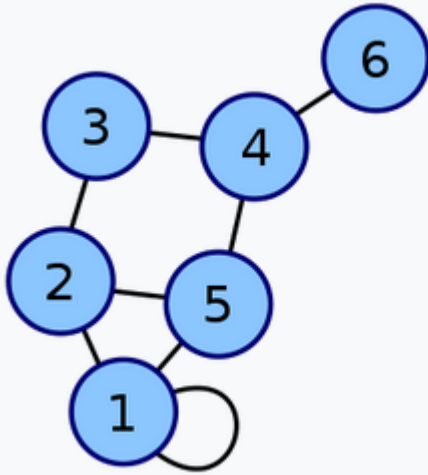


	$e_1$	$e_2$	$e_3$	$e_4$
<b>1</b>	1	1	1	0
<b>2</b>	1	0	0	0
<b>3</b>	0	1	0	1
<b>4</b>	0	0	1	1

$$= \begin{pmatrix} 1 & 1 & 1 & 0 \\ 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 1 \\ 0 & 0 & 1 & 1 \end{pmatrix}$$

[https://en.wikipedia.org/wiki/Incidence\\_matrix](https://en.wikipedia.org/wiki/Incidence_matrix)

# Adjacency Matrix

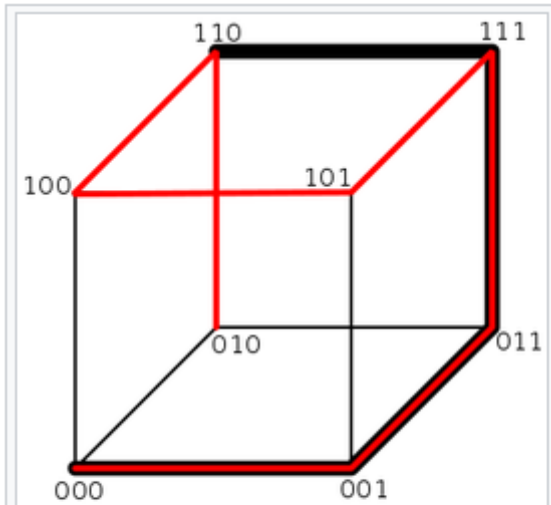


$$\begin{pmatrix} 2 & 1 & 0 & 0 & 1 & 0 \\ 1 & 0 & 1 & 0 & 1 & 0 \\ 0 & 1 & 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 & 1 & 1 \\ 1 & 1 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 \end{pmatrix}$$

Coordinates are 1-6.

[https://en.wikipedia.org/wiki/Adjacency\\_matrix](https://en.wikipedia.org/wiki/Adjacency_matrix)

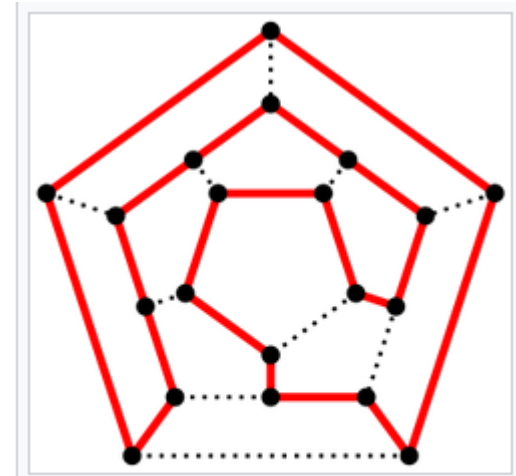
# Hamiltonian Path



A hypercube graph showing a Hamiltonian path in red, and a longest induced path in bold black.



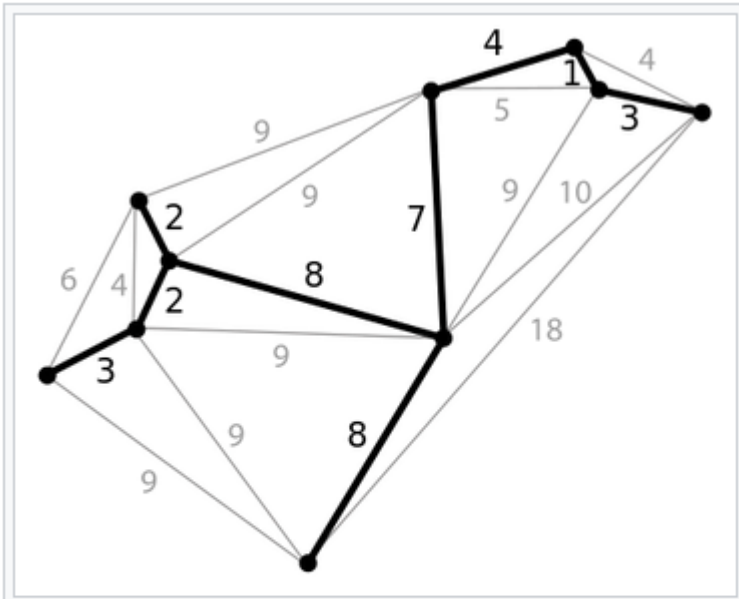
One possible Hamiltonian cycle through every vertex of a dodecahedron is shown in red - like all platonic solids, the dodecahedron is Hamiltonian



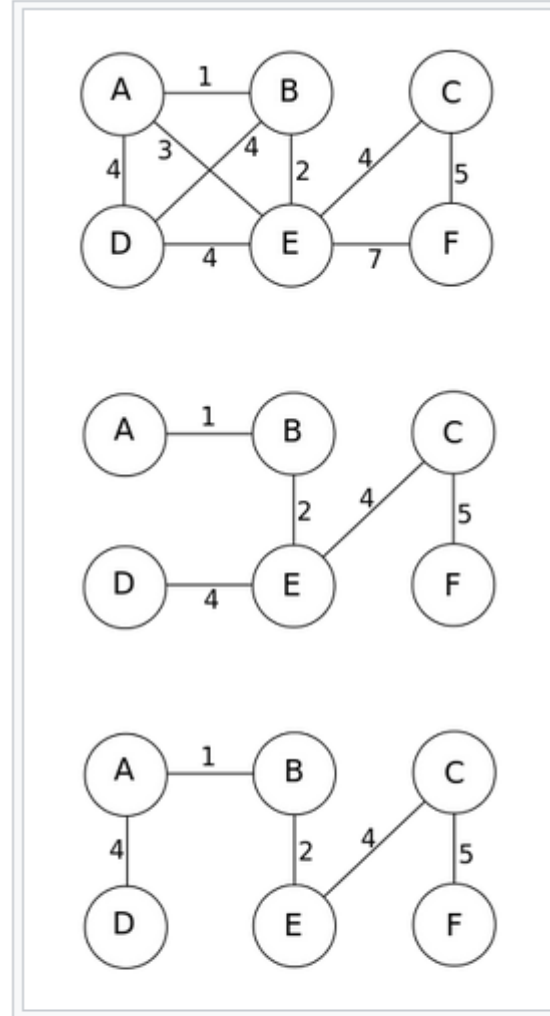
The above as a two-dimensional planar graph

[https://en.wikipedia.org/wiki/Path\\_\(graph\\_theory\)](https://en.wikipedia.org/wiki/Path_(graph_theory))

# Minimum Spanning Tree



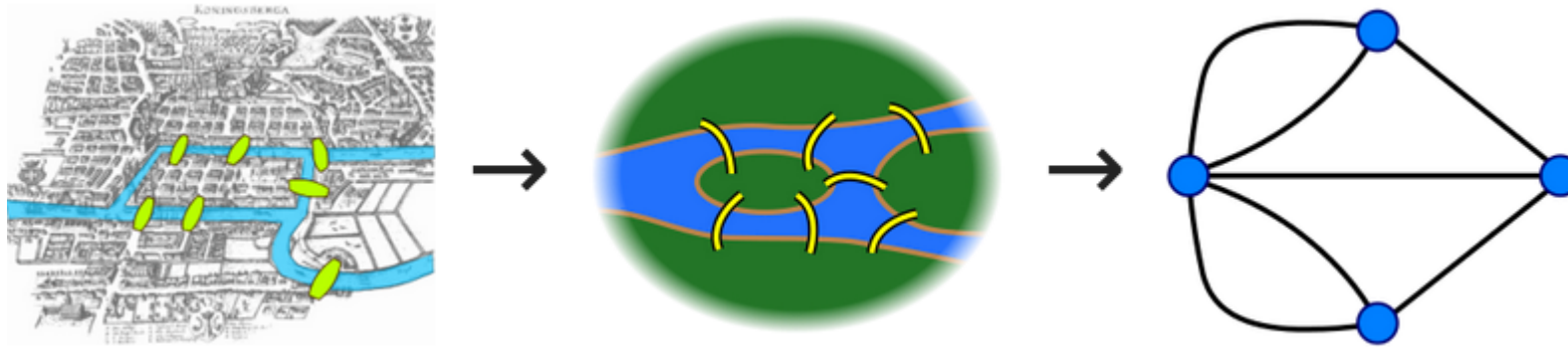
A **planar graph** and its minimum spanning tree. Each edge is labeled with its weight, which here is roughly proportional to its length.



This figure shows there may be more than one minimum spanning tree in a graph. In the figure, the two trees below the graph are two possibilities of minimum spanning tree of the given graph.

[https://en.wikipedia.org/wiki/Minimum\\_spanning\\_tree](https://en.wikipedia.org/wiki/Minimum_spanning_tree)

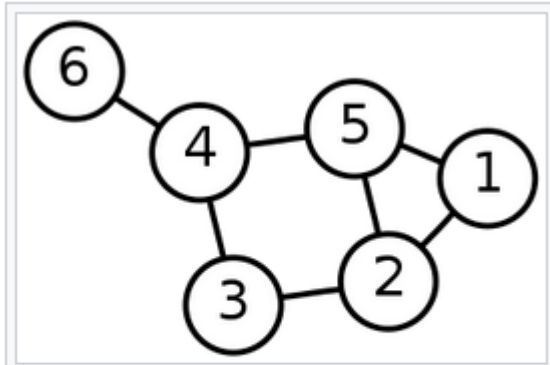
# Seven Bridges of Königsberg



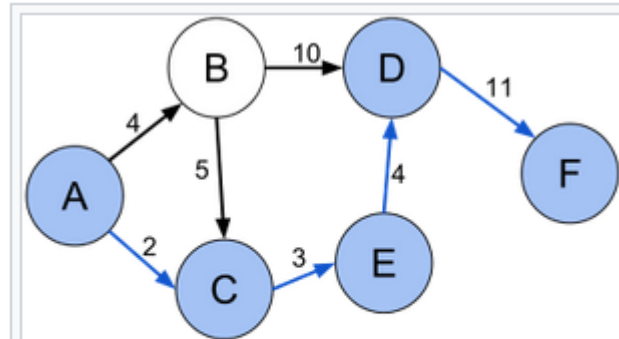
The problem was to devise a walk through the city that would cross each of those bridges once and only once.

[https://en.wikipedia.org/wiki/Seven\\_Bridges\\_of\\_K%C3%B6nigsberg](https://en.wikipedia.org/wiki/Seven_Bridges_of_K%C3%B6nigsberg)

# Shortest path problem



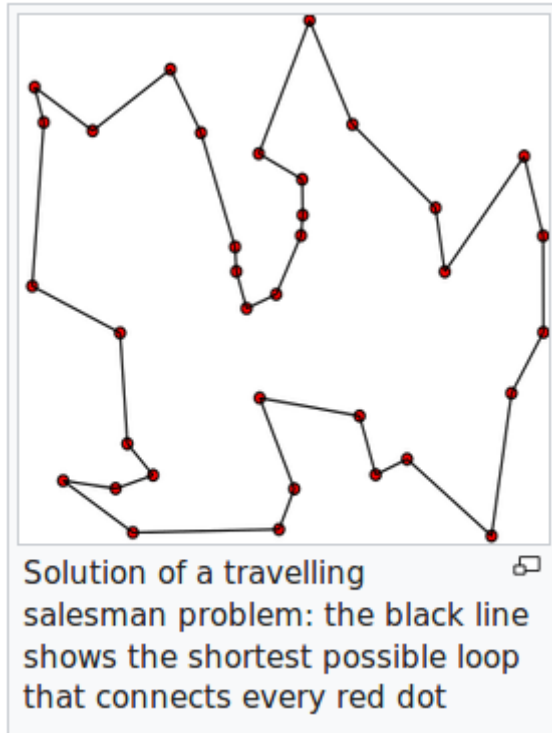
(6, 4, 5, 1) and (6, 4, 3, 2, 1) are both paths between vertices 6 and 1



Shortest path (A, C, E, D, F) between vertices A and F in the weighted directed graph

[https://en.wikipedia.org/wiki/Shortest\\_path\\_problem](https://en.wikipedia.org/wiki/Shortest_path_problem)

# Traveling salesman problem



[https://en.wikipedia.org/wiki/Travelling\\_salesman\\_problem](https://en.wikipedia.org/wiki/Travelling_salesman_problem)



# Simple Graph

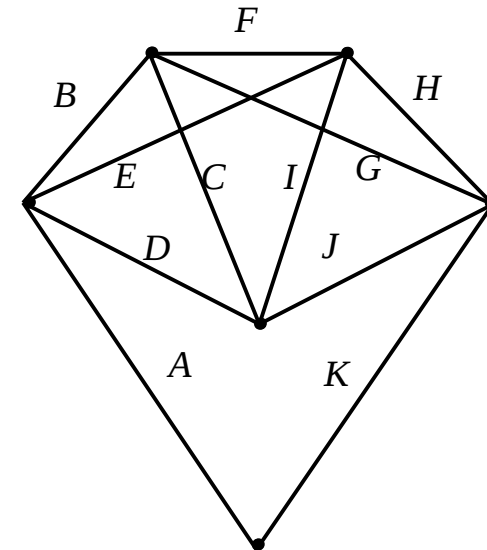
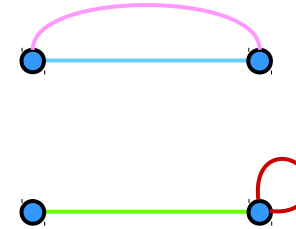
A simple graph is an undirected graph **without multiple edges or loops.**

the edges form a set (rather than a multiset)  
each edge is an unordered pair of distinct vertices.

can define a simple graph to be a **set  $V$**  of vertices  
together with a **set  $E$**  of edges,

**$E$**  are 2-element subsets of  **$V$**

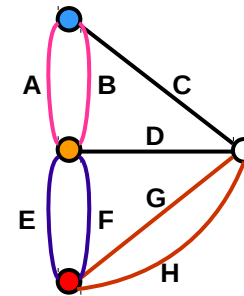
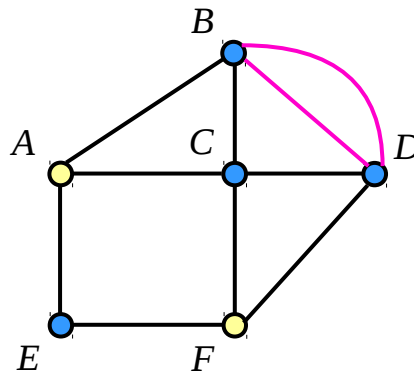
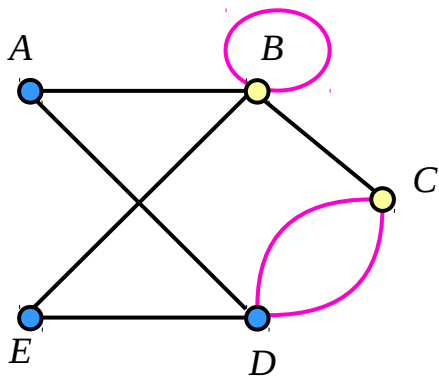
with  **$n$**  vertices,  
the **degree** of every vertex is at most  $n - 1$



[https://en.wikipedia.org/wiki/Travelling\\_salesman\\_problem](https://en.wikipedia.org/wiki/Travelling_salesman_problem)

# Multi-Graph

A **multigraph**, as opposed to a **simple graph**, is an undirected graph in which **multiple edges** (and sometimes **loops**) are allowed.



[https://en.wikipedia.org/wiki/Travelling\\_salesman\\_problem](https://en.wikipedia.org/wiki/Travelling_salesman_problem)

# Multiple Edges

- multiple edges
- parallel edges
- Multi-edges



are two or more edges  
that are incident to the same two vertices

A **simple graph** has no multiple edges.

[https://en.wikipedia.org/wiki/Travelling\\_salesman\\_problem](https://en.wikipedia.org/wiki/Travelling_salesman_problem)

# Loop

- a loop
- a self-loop
- a buckle

is an edge that connects a vertex to itself.

A **simple graph** contains no loops.



[https://en.wikipedia.org/wiki/Travelling\\_salesman\\_problem](https://en.wikipedia.org/wiki/Travelling_salesman_problem)

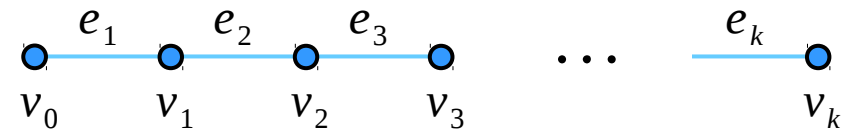
# Walks

For a graph  $G = (V, E)$ , a **walk** is defined as a sequence of alternating vertices and **edges** such as  $v_0, e_1, v_1, e_2, \dots, e_k, v_k$

where each edge  $e_i = \{v_{i-1}, v_i\}$

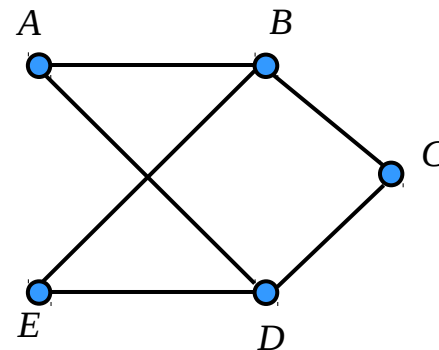
The length of this walk is  $k$

**Edges** are allowed to be repeated



$e_i = e_j$  for some  $i, j$

*ABCDE*  
*ABCDCBE*

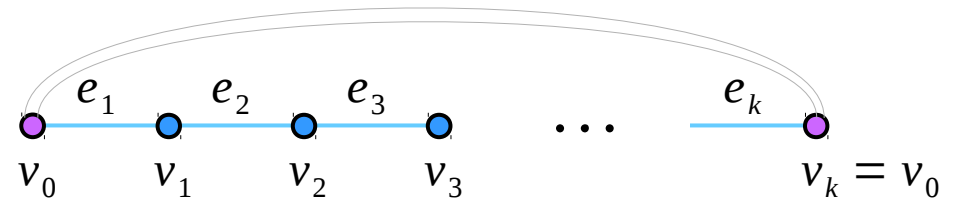


<http://mathonline.wikidot.com/walks-trails-paths-cycles-and-circuits>

# Open / Closed Walks

A walk is considered to be **closed** if the **starting** vertex is the same as the **ending** vertex.

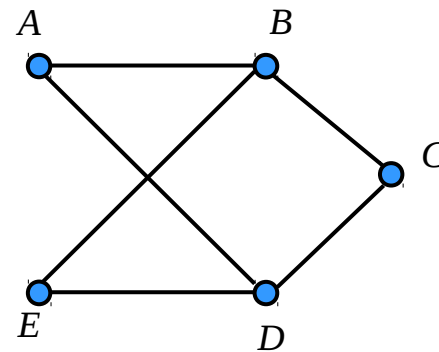
Otherwise **open**



*ABCD*A

*ABCDE*

*ABCD**C**B**E*

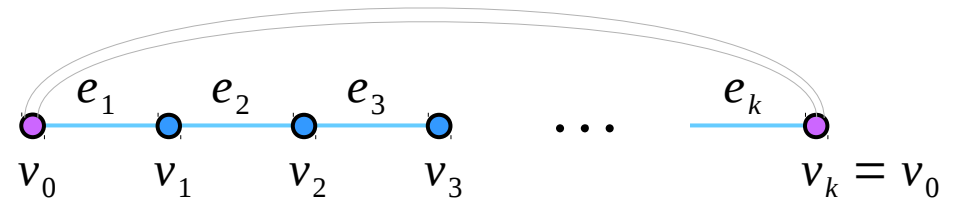


<http://mathonline.wikidot.com/walks-trails-paths-cycles-and-circuits>

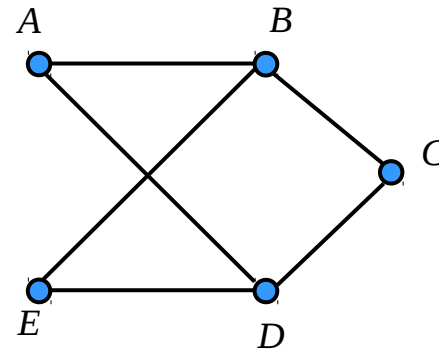
# Open / Closed Walks

A walk is considered to be **closed** if the **starting** vertex is the same as the **ending** vertex.

Otherwise **open**



closed walk  $ABCD A$   
open walk  $ABCDE$   
open walk  $ABCDCBE$

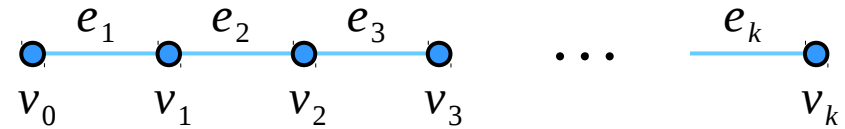


<http://mathonline.wikidot.com/walks-trails-paths-cycles-and-circuits>

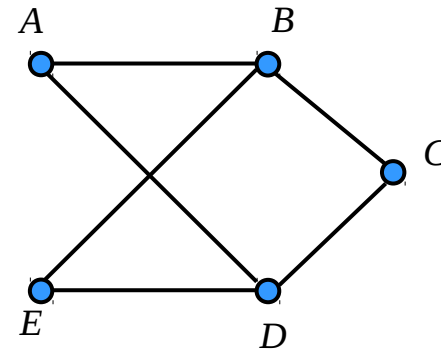
# Trails

A **trail** is defined as a **walk** with no repeated edges.

$$e_i \neq e_j \text{ for all } i, j$$



closed trail	closed walk	$ABCDA$
open trail	open walk	$ABCDE$
<del>open trail</del>	<del>open walk</del>	<del><math>ABCDCBE</math></del>



<http://mathonline.wikidot.com/walks-trails-paths-cycles-and-circuits>

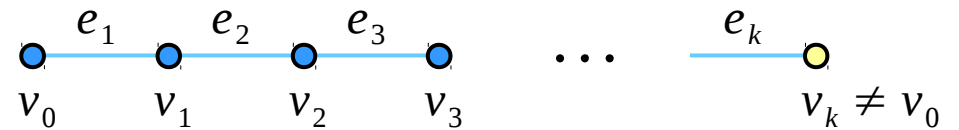


# Paths

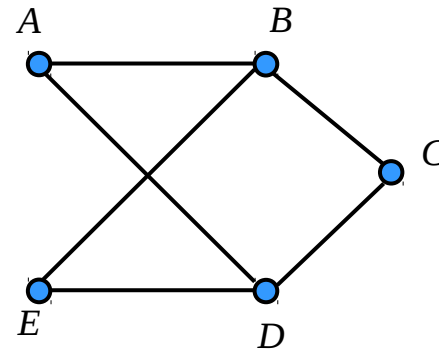
A **path** is defined as a **open trail** with no repeated vertices.

$$e_i \neq e_j \text{ for all } i, j$$

$$v_i \neq v_j \text{ for all } i, j$$



path	closed trail	closed walk	<i>ABCD</i>
path	open trail	open walk	<i>ABCDE</i>
path	open trail	open walk	<i>ABCD</i> <del><i>CB</i></del> <i>E</i>
path	open trail	open walk	<i>BED</i> <del><i>ABC</i></del>



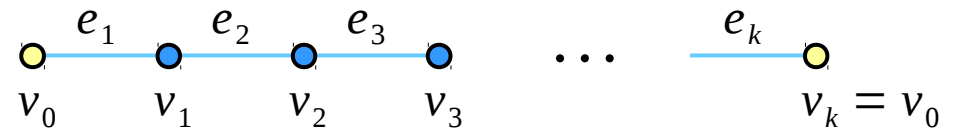
<http://mathonline.wikidot.com/walks-trails-paths-cycles-and-circuits>

# Cycles

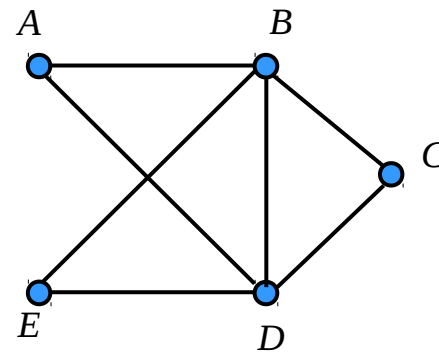
A **cycle** is defined as a **closed trail** with no repeated vertices except the **start/end vertex**

$$e_i \neq e_j \text{ for all } i, j$$

$$v_i \neq v_j \text{ for all } i, j$$



cycle	circuit	closed walk	<i>ABCD</i>
eyele	circuit	closed walk	<i>ABCDEBDA</i>

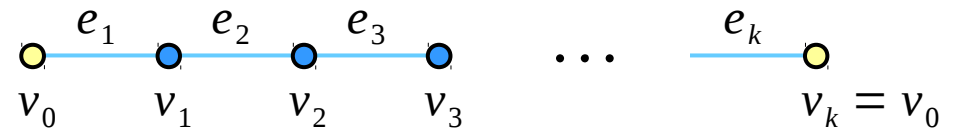


<http://mathonline.wikidot.com/walks-trails-paths-cycles-and-circuits>

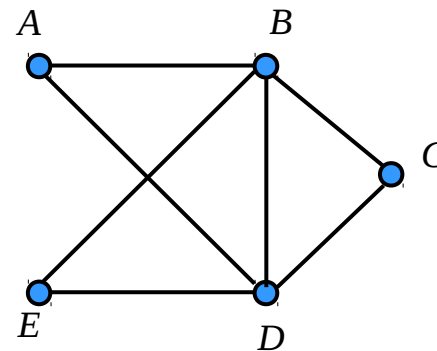
# Circuits

A **circuit** is defined as a **closed trail** with possibly repeated vertices but with no repeated edges

$$e_i \neq e_j \text{ for all } i, j$$
$$v_i = v_j \text{ for some } i, j$$

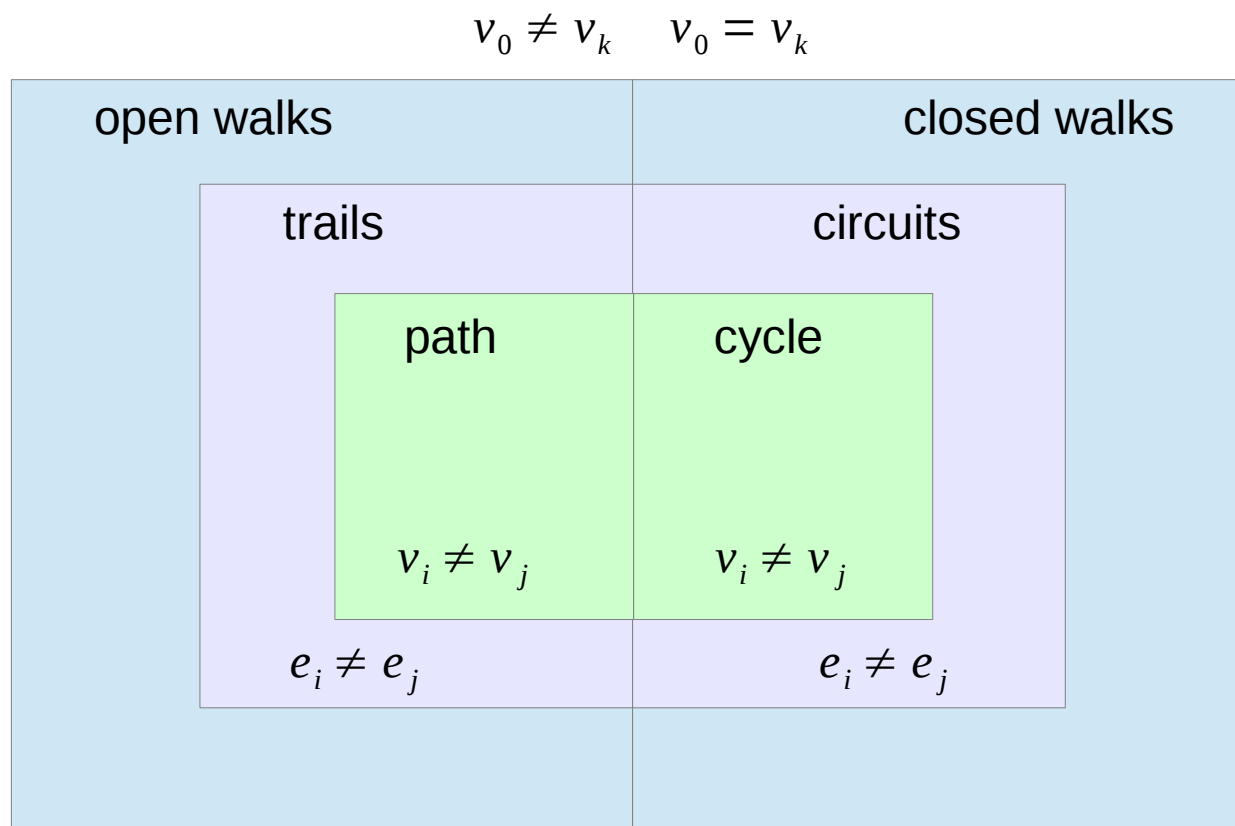


circuit	closed walk	$ABCD A$
circuit	closed walk	$ABCDEBDA$

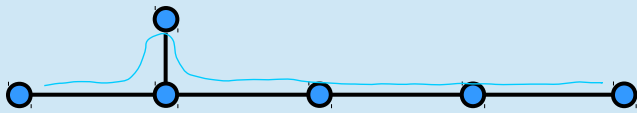
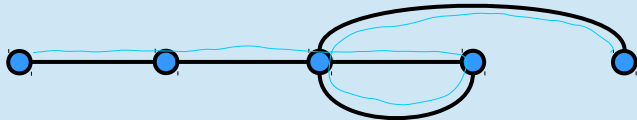

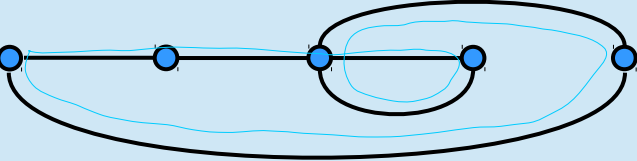
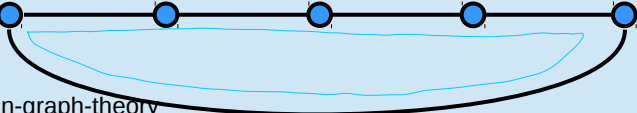


<http://mathonline.wikidot.com/walks-trails-paths-cycles-and-circuits>

# Walk, Trail, Path, Circuit, Cycle



# Walk, Trail, Path, Circuit, Cycle

	<b>Vertices</b>	<b>Edges</b>		
<b>Walk</b>	may repeat	may repeat	(Closed/Open)	
<b>Trail</b>	may repeat	<u>cannot</u> repeat	(Open)	
<b>Path</b>	<u>cannot</u> repeat	<u>cannot</u> repeat	(Open)	
<b>Circuit</b>	may repeat	<u>cannot</u> repeat	(Closed)	
<b>Cycle</b>	<u>cannot</u> repeat	<u>cannot</u> repeat	(Closed)	

<https://math.stackexchange.com/questions/655589/what-is-difference-between-cycle-path-and-circuit-in-graph-theory>

## References

- [1] <http://en.wikipedia.org/>
- [2]